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# Method of Recovery of Deleted Records in a Postgre SQL Database

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#### METHOD OF RECOVERY OF DELETED RECORDS IN A POSTGRESQL DATABASE

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Keywords:
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Hex editor

Received: 28 April 2017 Accepted: 04 June 2017 Published: 17 August 2017 **Abstract.** This paper recommends recovery methods for records deleted from PostgreSQL. This paper includes a description of PostgreSQL, which is the subject of this study. It also incorporates a layout of the pages and records relevant to data storage and the deletion tests and recovery algorithm used. When recovering deleted records, PostgreSQL data files, which contain a history of deleted records, are required. Also, a Hex editor is used, which can check the contents of data files. One can analyze each record's header information by opening PostgreSQL data files with a Hex editor. Then, the analyzed header information from the records is checked. If it is deemed that there are deleted records, they can be extracted for recovery. If VACUUM, a record cleanup program offered by PostgreSQL, has been used for deletion, the records cannot be restored. In this study, it was possible to recover the deleted records of the PostgreSQL database in the damages caused by cybercrime. Recovered records can be used as legal evidence, which can benefit companies in a court situation. The benefit of being in a court situation helps protect the interests of the enterprise.

#### INTRODUCTION

As enterprises depend significantly on IT systems, and as the amount of handled information increases, the use of databases grows and their range of usage also widens. With the increased use of databases, and a rise in IT-related corporate crimes, more and more database records are being permitted as legal evidence. However, database records can be deliberately or routinely deleted. Therefore, in order for database records to be allowed as legal evidence, these deleted records should be recovered.

Previous studies on the methods used to restore records deleted from databases have only been carried out in regard to Oracle and certain DBMS. Research on the recovery methods for records deleted from PostgreSQL, an open source database, has not yet been carried out. PostgreSQL is a very important database with 4th largest market share in the world, but there is no mention of PostgreSQL in other studies and no research has been done [5]. We have studied PostgreSQL's record recovery method, which was not existing. Hence, this paper recommends methods of recovery for records deleted in PostgreSQL. With respect to the methods used to recover deleted records, PostgreSQL data files, which feature a history of deleted records, are necessary. Also, a Hex editor, which checks the content of the data files, is used. PostgreSQL data files are opened with the Hex editor, and then each record's header information is analyzed. If it is deemed that the record was deleted, it is extracted and recovered.

This paper analyzes database changes using generation and deletion tests for records; it also suggests recovery methods for deleted records based on the analysis, identifies exceptional cases, and investigates records' retrievability.

This paper includes a description of PostgreSQL, which is the subject of this study, and it also incorporates a layout of the pages and records relevant to data storage, and the deletion tests and recovery algorithm that were used.

#### LITERATURE REVIEW

DBMS, whose recovery methods for deleted records were studied, include Oracle, DB2, and the SQL Server, among other commercial DBMS, as well as MySQL (an open-source DBMS) [1], [2], [3], [4]. As for the relevant research, we will select two DBMSs among the commercial DBMSs, which have the top two highest market shares, and we will investigate their recovery methods.

In this paper, we will investigate the recovery methods for the records and tables in PostgreSQL. In addition, we will identify the threshold at which records and tables cannot be restored, and we will also discuss any pertinent test results.

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#### The Method of Recovery for Deleted Records in the Oracle Database

To restore records deleted from Oracle, the System tablespace and user tablespace created by the user are required. By figuring out the blocks of user tablespace, the record contents of every table can be identified, and one can also determine whether the records were deleted according to the flag value for that record. When deleting records, one flag value for each record of the data block changes. Flag values change from 0x2C to 0x3C, while other information remains fixed [1].

When dropping a table, two flag values for the deleted tables of the OBJ\$ schema and C\_OBJ\$ schema in system tablespace change. The flag value for the deleted tables in the OBJ\$ schema changes from 0x2C to 0x3C, and the flag value for the deleted tables in the C-OBJ\$ schema changes from 0x6C to 0x7C.

Exceptionally, when the Shrink feature, which cleans up deleted space to optimize the database, is used, the deleted records and tables cannot be restored. Recovery of the tables is only possible when the records are not overwritten.

After extracting the deleted records and saving them as separate files, the restoration of deleted records ends. Table recovery was not mentioned in the relevant research.

## The Method of Recovery for Deleted Records in DB2 Database

When deleting records from DB2, the flag value for EX-TENT MAP inside the tablespace changes. When generating records, the flag values are given sequential values, and when they are deleted, the flag value changes to 0xFFFF while the records' other information remains fixed.

When dropping a table, the records are kept and they can be restored. However, the flag values that are related to deleting a table are not identified in the relevant research [2].

There are exceptional cases in which records cannot be recovered; this is when the last records are deleted and when new records are generated. In this case, it is impossible to recover the record, for the information on the last deleted record has been overwritten. In addition, when the database inside the DBMS is deleted with the DROP command, the records included in the deleted database cannot be recovered.

After extracting the deleted records and saving them as separate files, restoration work for the deleted records ends. There is no reference to the restoration work required for tables in the relevant research.

#### METHOD AND MATERIALS PostgreSQL

PostgreSQL is an Object-Relational Database Management System (ORDBMS) that is based on POSTGRES 4.2, which was developed by UC Berkeley's Computer Science Department [6].

PostgreSQL supports representative functionalities, such as complex syntax, foreign key, trigger, updatable view, transaction integrity, Multi-Version Concurrency Control (MVCC), etc. and it follows SQL standards. Also, PostgreSQL can be extended by users in a number of ways. Functionalities that can be extended include data type, function, operator, aggregate function, index method, procedural language, etc.

When deleting records in PostgreSQL, the DELETE command is used. The format of the DELETE command is "DELETE FROM table name WHERE column name = 'column value'". When not using the WHERE condition and using the DELETE FROM table name together, all records in the table are deleted. Since there is no additional confirmation for deletion when deleting records, one must be careful. This research was carried out based on PostgreSQL 9.4.

#### **Data Directory and File Structure**

PostgreSQL defines the settings and data files used in DBMS as PGDATA. Generally, PGDATA means /var/lib/pgsql/data. PGDATA includes several subdirectory and control files, which are shown in TABLE 1. Also, settings files such as postgresql.conf, pg\_hba.conf, and pg\_ident.conf are included in PGDATA, but they can be saved in other places as well [7], [8], [11].



	DAMADINECTORY MAD TIELS
Item	Description
postmaster.pid	A lock file written the current process ID, data directory path, start
timestamp, port,	unix domain socket path, listen address, shared memory segment ID
postgresql.auto.conf	A file for storing config parameters that are set by alter system
postmaster.opts	A file for command line options the server was started with
PG_VERSION	A file written major version information
base	Directory including data directory
global	Directory including system table
pg_clog	Directory including status of transaction
pg_dynshmem	Directory including files that used by dynamic shared memory's subsystem
pg_logical	Directory including status data for logical decoding
pg_multixact	Directory including data of multi transaction status
pg_notify	Directory including data of listen/notify status
pg_replslot	Directory including data of replication slot
pg_serial	Directory including information about committed serializable transactions
pg_snapshots	Directory including experted snapshot
pg_stat	Directory including permanent file about statistics
pg_stat_tmp	Directory including temporary file for statistics
pg_subtrans	Directory including data of subtransaction status
pg_tblspc	Directory including symbolic link for tablespace
pg_twophase	Directory including status of prepared transaction
pg_xlog	Directory including write ahead log file

TABLE 1 DATA DIRECTORY AND FILES

The database generated by a user is generated as a subdirectory 'base', which is a subdirectory of PGDATA. The directory name of the database generated by the user becomes OID (ObjectID), which is automatically generated in the DBMS. The table name generated by the user becomes OID, which is automatically generated in the DBMS. Temporary tables are generated in tAAA\_BBB form; AAA serves as identification of the process that generated the file, and BBB becomes an OID that is automatically generated in the DBMS.

#### Page Layout

Every table is stored as an array of pages of a fixed size. The size of the page is 8 KB. In a table, since all pages are logically equivalent, the record can be saved in any page. Each page consists of 5 parts, and the layouts comprising the page are shown in TABLE 2 [9], [12].

TABLE 2 PAGE LAYOUT

Item	Description
Page Header	Include general information about page, and also had an unallocated space information. total 24 bytes.
Record identifier	Information about actual records location. 4 bytes per record.
Unallocated space	New record identifier is allocated from the start of this area, and actual new record is allocated from the end.
Record	The actual record itself.
Special space	For special space for specific data(optional)



2017

PAGE HEADER LAYOUT Flag Length Description pd\_lsn 8 bytes About next byte after last byte of xlog record for last change to this page pd\_checksum 2 bytes Information of page checksum pd\_flags 2 bytes Information of flag bits pd\_lower 2 bytes About location. Offset for start of unallocated space pd\_upper 2 bytes About location. Offset for end of unallocated space pd\_special 2 bytes About location. Offset for start of special space pd\_pagesize\_version 2 bytes Information about page size and layout version For oldest unpruned XMAX on page. pd\_prune\_xid 4 bytes

The first 24 bytes of each page are the page header. The

layouts comprising the page header are shown in TABLE 3.

TABLE 3

The record identifiers are followed by the page header. Record identifiers require 4 bytes each. Record identifiers contain the start of the record, size, and certain attributes of the record. New record identifiers are allocated at the start of the unallocated space.

Unallocated space is followed by a record identifier. The unallocated space is where information is not stored.

The records are stored at the end of the unallocated space. The records themselves are where the contents entered by the user are stored.

The final section is the special section, which is generally not used.

#### **Record Layout**

All record structures are the same, and they contain fixed header sizes and other optional flag. The optional parts are t\_infomask, t\_hoff, and data entered by the user. Layouts comprising the header are shown in Table 4. The starting point of the data stored by the user is a byte that comes after the t\_infomask2 value.

	RE	CORD HEADER LAYOUT
Flag	Length	Description
t_xmin	4 bytes	TransactionID stamp for insert
t_xmax	4 bytes	TransactionID stamp for delete
t_cid	4 bytes	CommandID stamp for insert and(or) delete
t_xvac	4 bytes	TransactionID for vacuum operation
t_ctid	6 bytes	current TID of this or newer row
t_infomask2	2 bytes	attributes number and various flag bits
t_infomask	2 bytes	various flag bits(optional)
t_hoff	1 byte	user data's offset(optional)

TABLE 4

#### Change of Data file Triggered by Generating Record

There are four changes in page layout when record is generated. First, information of page header layout is changed. Second, record identifier is changed. Third, record data are added. Fourth, unallocated space is shrunken. Fig. 1 is figured about before the generating record.



Offset(h)	00	01	02	03	04	05	06	07	08	09	0A	0B	0C	0D	0E	OF	page header
00000000	00	00	00	00	B8	C4	7D	01	00	00	00	00	1C	00	DO	1F	,Ä}Ð.
00000010	00	20	04	20	00	00	00	00	DO	9F	60	00	00	00	00	00	П`
00000020	00	00	00	00	00	00	00	00	00	00	00	00	00	00	-00.	00	
																	record-1 identifier
						una	allo	ate	d s	pace	e						
																	record-1
00001 FA0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	/
00001FB0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
00001FC0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	/
00001FD0	49	07	00	00	00	00	00	00	00	00	00	00	00	00	00	00	I
00001FE0	01	00	03	00	02	08	18	00	6F	00	00	00	1B	61	61	61	aaa
00001FF0	61	61	61	61	20	20	20	20	20	00	00	00	E2	16	00	00	aaaaâ



Fig. 2 is figured about after the generating record.



Fig. 2. Page layout after generating record



#### Change of Data file Triggered by Record Deletion

There are two changes in record header layout when record is deleted. First, t\_xmax flag is changed. Before deletion, the value of t\_xmax is '0x00 00 00'; afterwards, an unspecified value is allocated. Second, t\_ctid flag is changed. According to the deletion test, the value of 4th byte from the

t\_ctid value changes from '0x00' to '0x20'; the value of 6th byte from the t\_ctid value changes from unspecified value to unspecified value. Only, the value of 4th byte from the t\_ctid value is changed regularly. Fig. 3 is figured about after the deleting record.

t_d	tid 4th value	/ t_xmax	/t_ctid 6th va	lue	
00001FA0	4A 07 00 00	00 00 00 00	0 00 00 00 00	00 00 00 00	J
00001FB0	02 00 03 00	02 08 18 0	0 DE 00 00 00	1B 62 62 62	Þbbb
00001FC0	62 62 62 62	20 20 20 2	0 20 00 00 00	01 17 00 00	bbbb
	ł	After variou	sly test, "t_ctid 4	4th value" is cha	nged regularly
00001FA0	4A 07 00 00	4B 07 00 0	0 00 00 00 00	00 00 00 00	JK
00001FB0	02 00 03 20	02 05 18 0	0 DE 00 00 00	1B 62 62 62	Þbbb
00001FC0	62 62 62 62	20 20 20 2	0 20 00 00 00	01 17 00 00	bbbb

Fig. 3. After the deleting record

Instances, where the record header information changed according to performing the deletion test, were above three results. For all continuous tests, the 4th byte from the  $t_{\rm L}$  tid value after the change remained fixed.

When distinguishing the deleted records, if the 4th byte from the t\_ctid value was '0x00', it was not deleted, and if it was '0x20', it was found to be deleted.

**Record Generation of after Deletion** 

When a new record is generated after the existing one is deleted, the contents entered by the user of the deleted record are not overwritten, and they remain without any changes. Deletion of the existing record and generation of a new record are shown in Fig. 4.

00001FA0	4A	07	00	00	4B	07	00	00	00	00	00	00	00	00	00	00	JK
00001FB0	02	00	03	20	02	05	18	00	DE	00	00	00	1B	62	62	62	Þbbb
00001FC0	62	62	62	62	20	20	20	20	20	00	00	00	01	17	00	00	bbbb
00001 FD0	49	07	00	00	00	00	00	00	00	00	00	00	00	00	00	00	т
00001250	01	00	0.2	00	02	00	10	00	6F	00	00	00	18	61	61	61	
00001450	UI	00	05	00	02	03	10	00	OF	00	00	00	TD	OT	OT	OT	
00001FF0	61	61	61	61	20	20	20	20	20	00	00	00	E2	16	00	00	aaaaâ
/			1														
deleted rec	ord	1															
	+																
exis	t ree	coru															
			200														
		-	ne	w re	ecor	d											
00001F10	4C	07	00	w re	ecor	<b>d</b>	00	00	00	00	00	00	00	00	00	00	L
00001F10 00001F80	4C 03	07	00 03	00 00	00 02	d 00 09	00	00	00 4D	00	00	00	00 1B	00	00	00 63	LMccc
00001F10 00001F80 00001F90	4C 03 63	07 00 63	00 03 63	00 00 63	00 02 20	00 09 20	00 18 20	00 00 20	00 4D 20	00 01 00	00000	00 00 00	00 1B 1E	00 63 17	00 63 00	00 63 00	LMccc cccc
00001F10 00001F80 00001F90 00001FA0	4C 03 63 4A	07 00 63 07	00 03 63 00	00 00 63 00	00 02 20 4B	00 09 20 07	00 18 20	00 00 20 00	00 4D 20	00 01 00	00 00 00	00 00 00	00 1B 1E 00	00 63 17 00	00 63 00	00 63 00	LMccc cccc JK.
00001F10 00001F80 00001F90 00001FA0 00001FB0	4C 03 63 4A 02	07 00 63 07 00	00 03 63 00 03	00 00 63 00 20	00 02 20 4B 02	00 09 20 07 05	00 18 20 00 18	00 00 20 00	00 4D 20 00 DE	00 01 00 00	00 00 00 00	00 00 00 00	00 1B 1E 00 1B	00 63 17 00 62	00 63 00 00 62	00 63 00 00 62	LMccc cccc JK
00001F10 00001F80 00001F90 00001FA0 00001FB0 00001FC0	4C 03 63 4A 02 62	07 00 63 07 00 62	00 03 63 00 03 62	00 00 63 00 20 62	00 02 20 4B 02 20	d 00 09 20 07 05 20	00 18 20 00 18 20	00 00 20 00 00 20	00 4D 20 00 DE 20	00 01 00 00 00	00 00 00 00 00 00	00 00 00 00 00	00 1B 1E 00 1B 01	00 63 17 00 62 17	00 63 00 00 62 00	00 63 00 00 62 00	LMccc cccc JK
00001F10 00001F80 00001F90 00001FA0 00001FB0 00001FC0 00001FD0	4C 03 63 4A 02 62 49	07 00 63 07 00 62 07	00 03 63 00 03 62 00	00 00 63 00 20 62 00	00 02 20 4B 02 20 00	00 09 20 07 05 20 00	00 18 20 00 18 20 00	00 00 20 00 20 20	00 4D 20 00 DE 20 00	00 01 00 00 00 00	00 00 00 00 00 00	00 00 00 00 00 00	00 1B 1E 00 1B 01 00	00 63 17 00 62 17 00	00 63 00 62 00 00	00 63 00 62 00 00	LMccc cccc JK bbbb L
00001F10 00001F80 00001F90 00001FA0 00001FB0 00001FC0 00001FD0 00001FE0	4C 03 63 4A 02 62 49 01	07 00 63 07 00 62 07 00	00 03 63 00 03 62 00 03	00 00 63 00 20 62 00 00	00 02 20 4B 02 20 00 00	d 00 09 20 07 05 20 00 00	00 18 20 00 18 20 00 18	00 20 00 20 20 20 00 00	00 4D 20 00 DE 20 00 6F	00 01 00 00 00 00 00	00 00 00 00 00 00 00	00 00 00 00 00 00 00	00 1B 1E 00 1B 01 00 1B	00 63 17 00 62 17 00 61	00 63 00 62 00 62 00 61	00 63 00 62 00 62 00 61	LMccc cccc JK bbbb bbbb Ioaaa

Fig. 4. Generated after the existing one is deleted



#### **Deletion of a Database and Table**

When deleting a database and table, the physical directory and the files targeted for deletion are deleted. Therefore, the records included in the targets for deletion cannot be restored. When deleting a table, the size of the physical file becomes 0 bytes; then, when restarting PostgreSQL, the physical file of the deleted table is deleted.

#### **Database Optimization Test**

Even if the records are deleted, a history of the records is kept. However, if VACUUM, which is a cleaning program that is offered as a basic program by PostgreSQL, is used, then the entire history of deleted records is also deleted [10].

#### RESULTS

When records are deleted, the value of t\_xmax and t\_ctid inside the record header changes. The value of t\_xmax changes

from '0x00 00 00 00' of 4 bytes to another value, which varies according to the value of the transaction. As for t\_ctid value, the value of fourth bytes among values comprised of 6 bytes changes from '0x00' to '0x20'.

#### **Recovery Algorithm for Deleted Records**

Copy the deleted table file to the system to be recovered. Connect to PostgreSQL and check the structure of the table to be restored. The table file to be restored is divided into pages and copied to a separate file. Each page copied into a separate file is divided into records. Check the t\_ctid flag of the record to separately collect records with the fourth byte value '0x20'. Combine the collected records with the table structure and convert the values of each column into ASCII form to complete recovery.

The deleted record recovery algorithm of PostgreSQL is shown in Fig. 5.



Fig. 5. Recovery algorithm for deleted records

#### DISCUSSION

The deleted record recovery method requires a PostgreSQL data file in which the structure of the table where the deleted record existed and the information of the deleted record exist [13,14]. Connect to PostgreSQL server, execute SQL command, and acquire table structure through it. To analyze the data file, copy the entire PostgreSQL data file to a computer that will perform the data recovery operation using a storage medium such as USB. The hex editor is used to check the contents of the data file with the records to be restored.

Open the PostgreSQL data file with a hex editor and check the header information of the record. If the value of the fourth byte of the t\_ctid flag of the confirmed header information of the deleted record is changed to '0x20', the record is determined as the deleted record. When the record is deleted, since the value of the fourth byte of the t\_ctid flag is changed consistently to '0x20', the t\_ctid flag is defined as a criterion for determining the deleted record.



The record information determined as the deleted record is stored as a separate file. By comparing the record information stored in a separate file with the structure of the table, it is determined which column value the record value to be restored is.

For the readability of the record information, the record information expressed in hexadecimal notation is converted into the ASCII form and the recovery operation of the deleted record is completed. If you use VACUUM, a cleanup program for deleted records provided by PostgreSQL, it is impossible to recover records because deleted record information is deleted.

#### CONCLUSION AND RECOMMENDATIONS

As IT usage in enterprises increases, the usage of databases increases accordingly. The increased use of databases raises the possibility that database records will need to be used as evidence in crime situations. For a database record to be utilized as legal evidence, records associated with the crime should be presented to the court. However, database records can be deleted both routinely and deliberately, which makes it impossible to present them to the court. Therefore, in order to submit deleted records to the court as evidence, it is necessary to restore the deleted records. A number of studies have been carried out on this topic, but they have only targeted record re-

covery methods for certain DBMSs. Studies of record recovery methods for PostgreSQL have not yet been performed, despite of the world's fourth largest market share. This paper proposes a record recovery method for PostgreSQL. We proposed a method to recover deleted records when PostgreSQL records are illegally deleted or deleted normally but need to be recovered. In this study, it was possible to recover the deleted records of the PostgreSQL database in the damages caused by cybercrime. Recovered records can be used as legal evidence. This can be beneficial to companies in court situations. The advantage of being in a court situation helps protect and keep the property of the enterprise. And the results of this study can be used to recover records even if normally deleted records are needed. However, there are limits to the possibility of recovering deleted records in some cases already mentioned.

Future work will be dedicated to studying for another DBMS that has not yet been studied. And we will also build a recovery automation solution if the environment is given.

#### **Declaration of Conflicting Interests**

This is the original work formulated and executed by the authors. It is not being processed or published elsewhere and there are no identifiable conflicts of interest.

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- This article does not have any appendix. -

